

The diagram illustrates a cryptographic algorithm structure, likely a Feistel network, enclosed in a dashed box labeled 124. The process begins with an **INPUT** block (64) labeled 110, which feeds into an **INITIAL PERMUTATION** block (112). The output of the initial permutation is split into two paths: **L0** (114) and **R0** (116). **L0** is labeled as the **PERMUTED INPUT** (122). Both **L0** and **R0** feed into a function block **f** (120), which also receives a key **K1** (118). The output of **f** is added to **L0** at a summation point (+). The result then splits into **L1 = R0** and **R1 = L0 + f(R0, K1)**. This structure repeats for subsequent rounds. Round **n** involves **L<sub>n</sub> = R<sub>n-1</sub>** and **R<sub>n</sub> = L<sub>n-1</sub> + f(R<sub>n-1</sub>, K<sub>n</sub>)**, where **K<sub>n</sub>** (156) is a round key. The final round shown is round **16**, where **L<sub>15</sub> = R<sub>14</sub>** and **R<sub>15</sub> = L<sub>14</sub> + f(R<sub>14</sub>, K<sub>15</sub>)**, followed by **R<sub>16</sub> = L<sub>15</sub> + f(R<sub>15</sub>, K<sub>16</sub>)** and **L<sub>16</sub> = R<sub>15</sub>**. The output of round 16 is labeled **PRE-OUTPUT** (126). This pre-output then passes through an **INVERSE INITIAL PERM.** block (130) to produce the final **OUTPUT** (64) labeled 132.

108

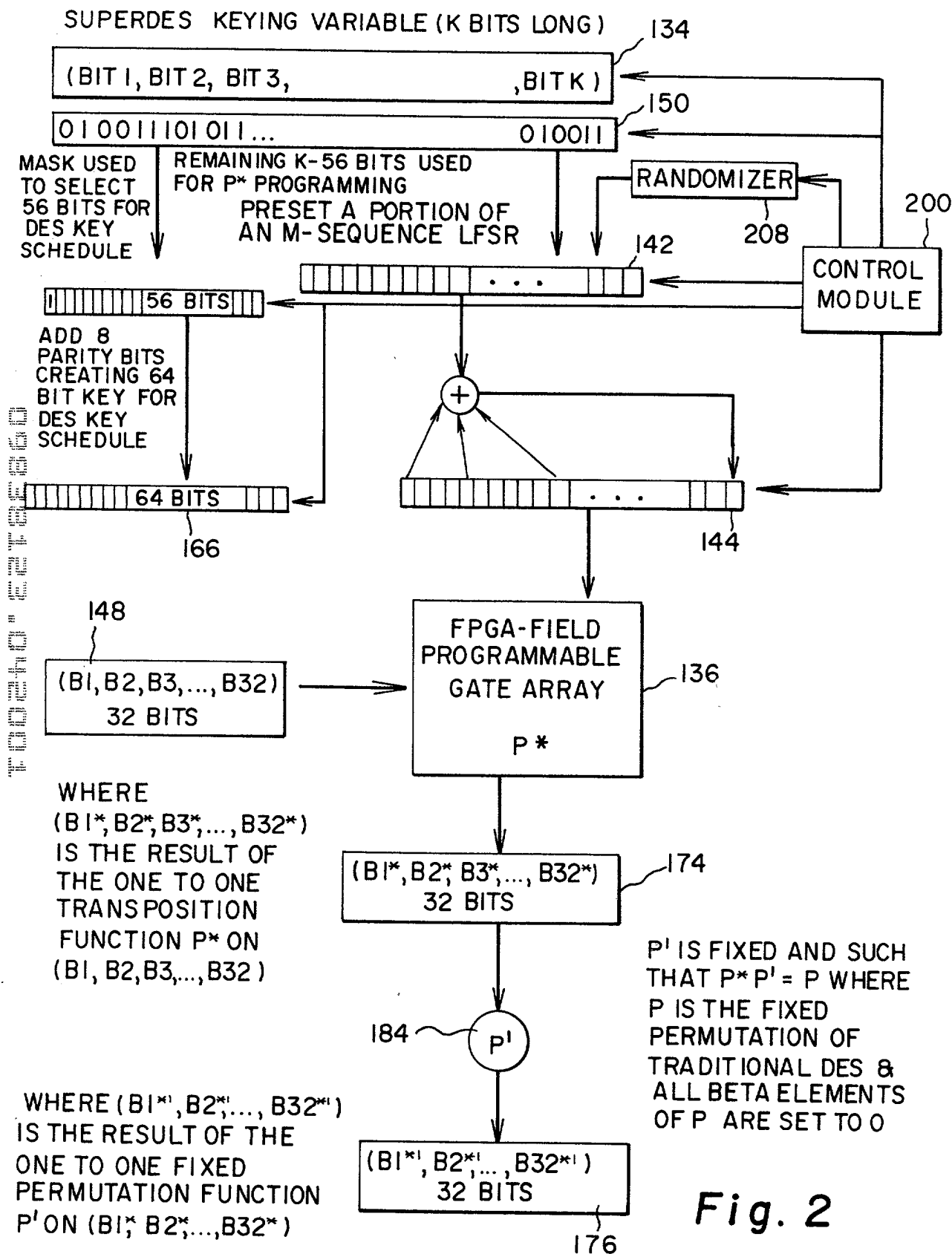


Fig. 2

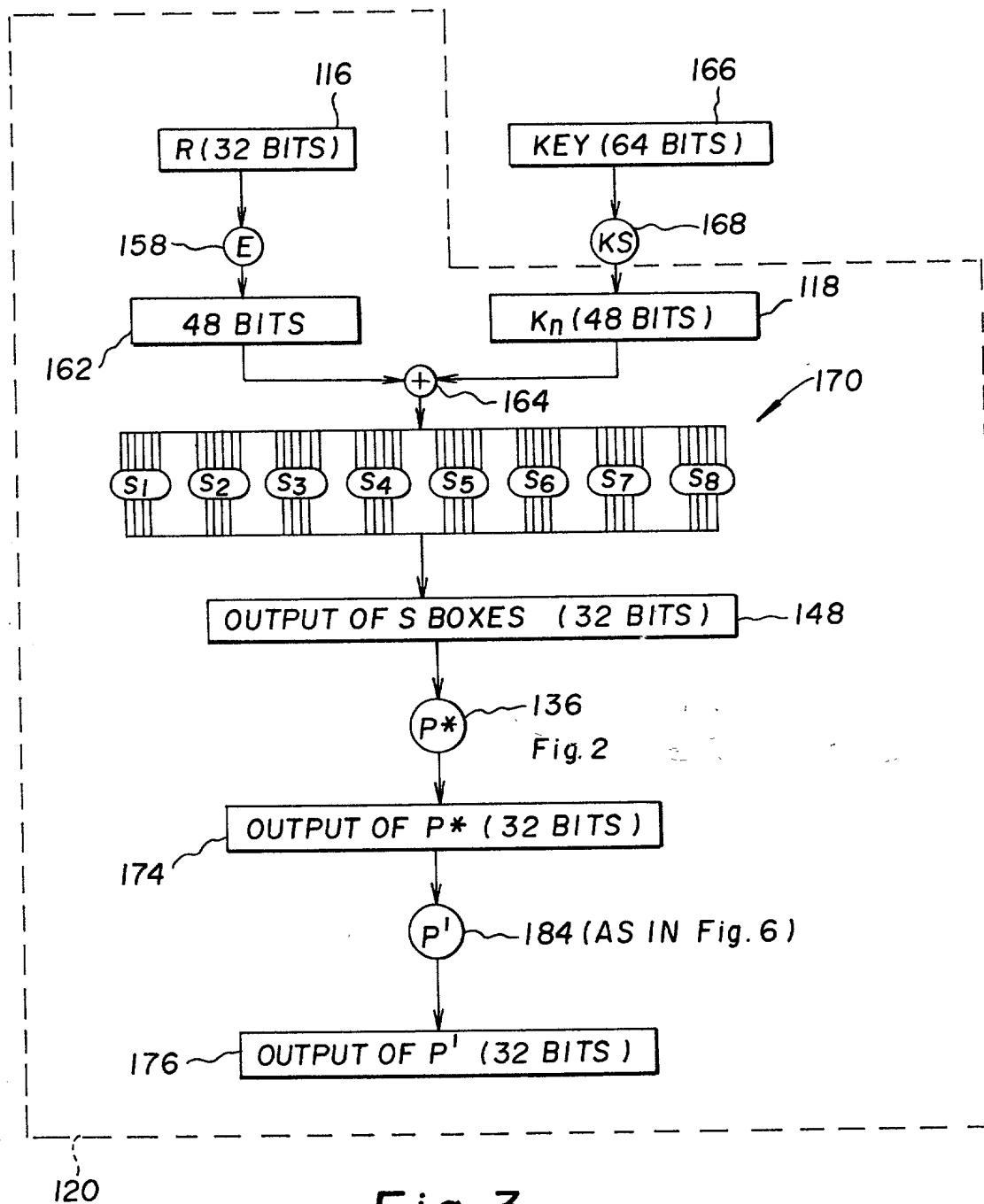
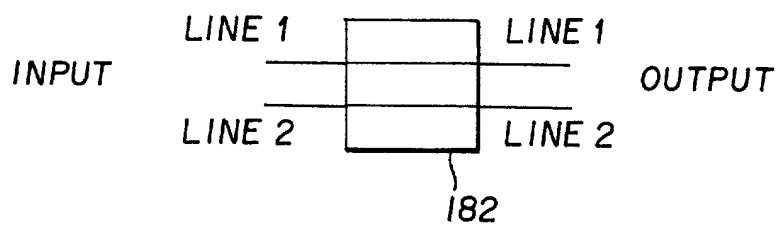
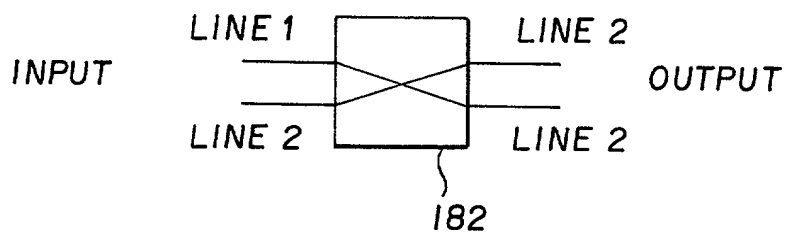


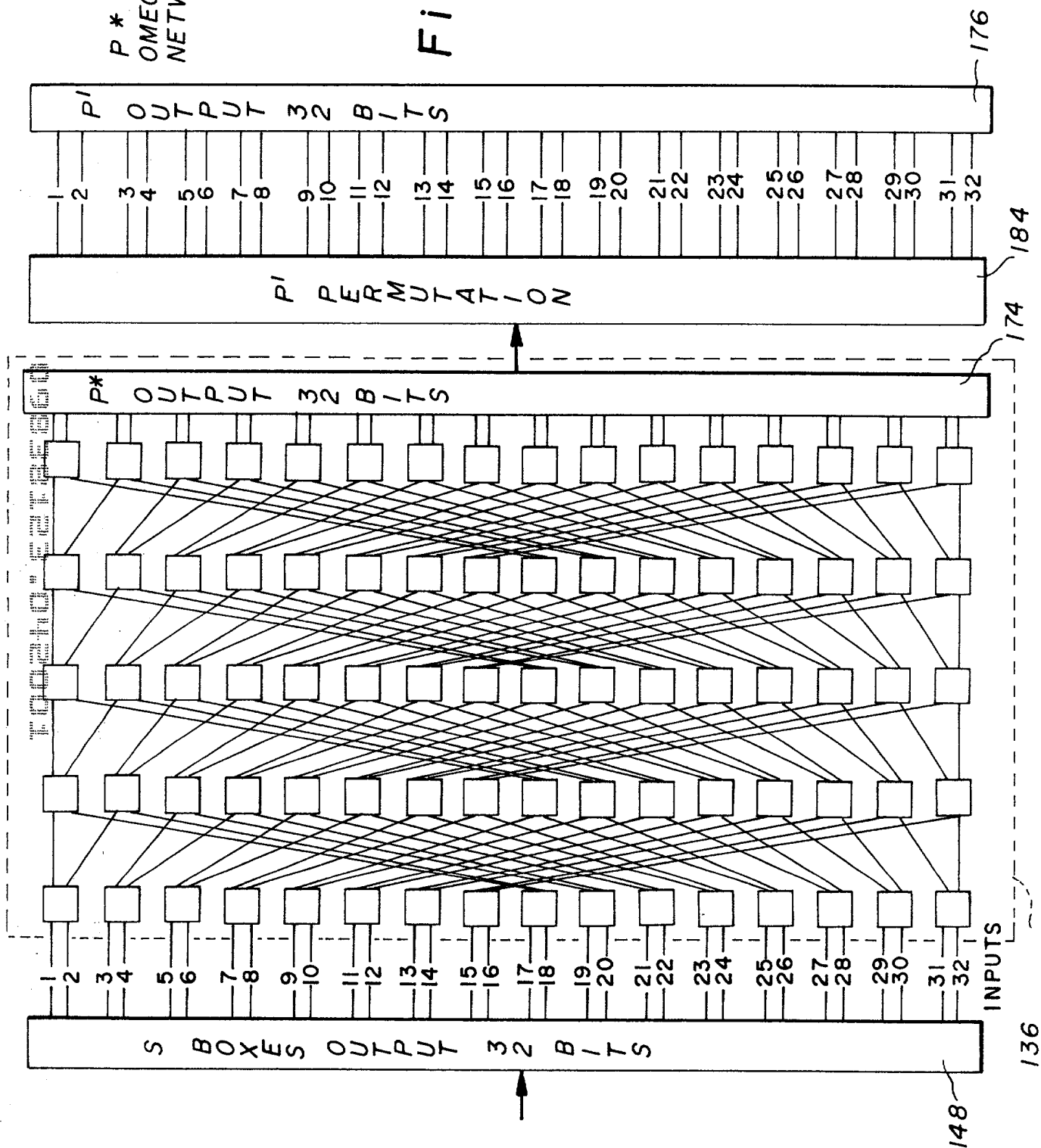
Fig. 3

$$\beta = 0$$

$$\beta = 1$$


**Fig. 5**

P \* IN AN  
OMEGA  
NETWORK

Fig. 6



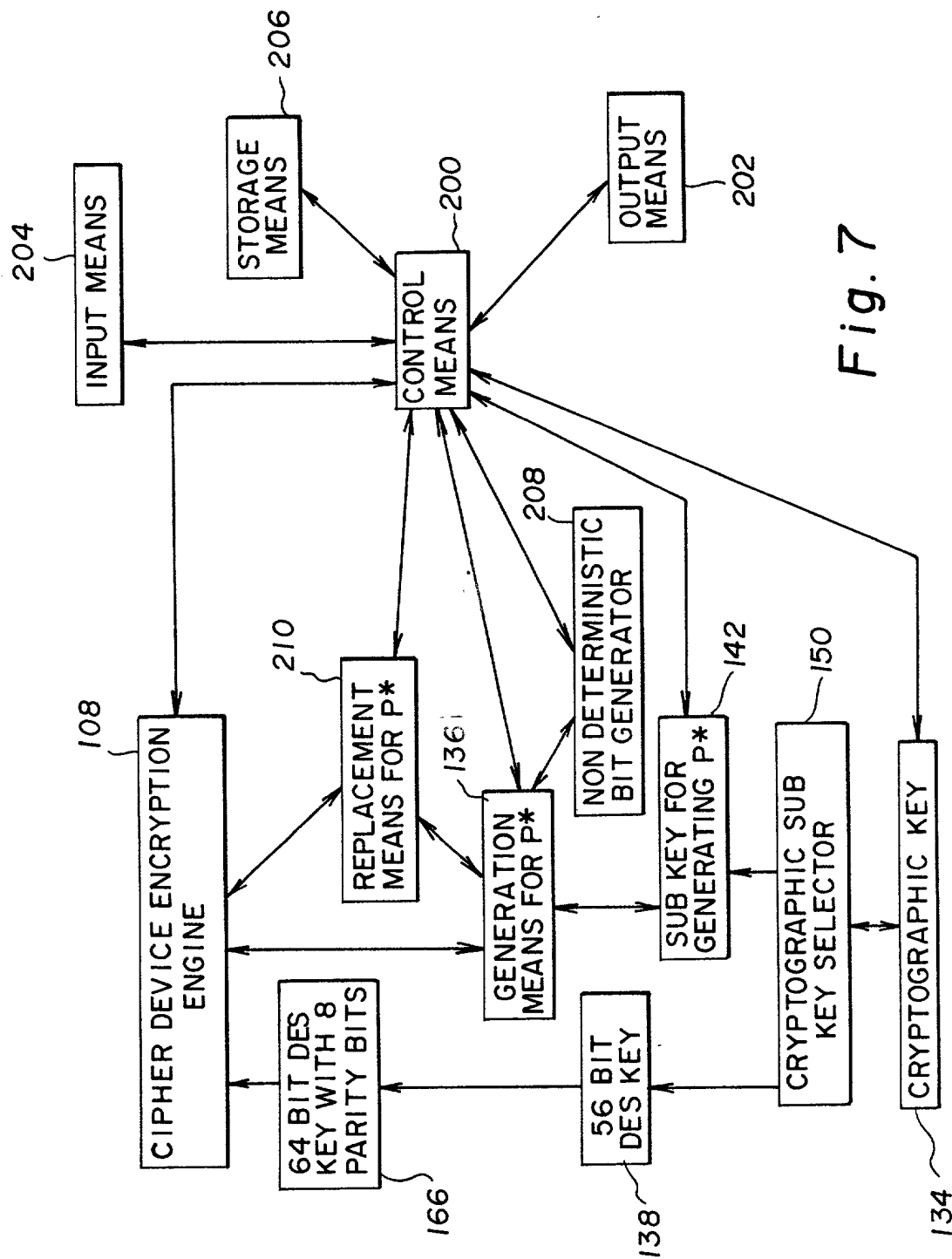
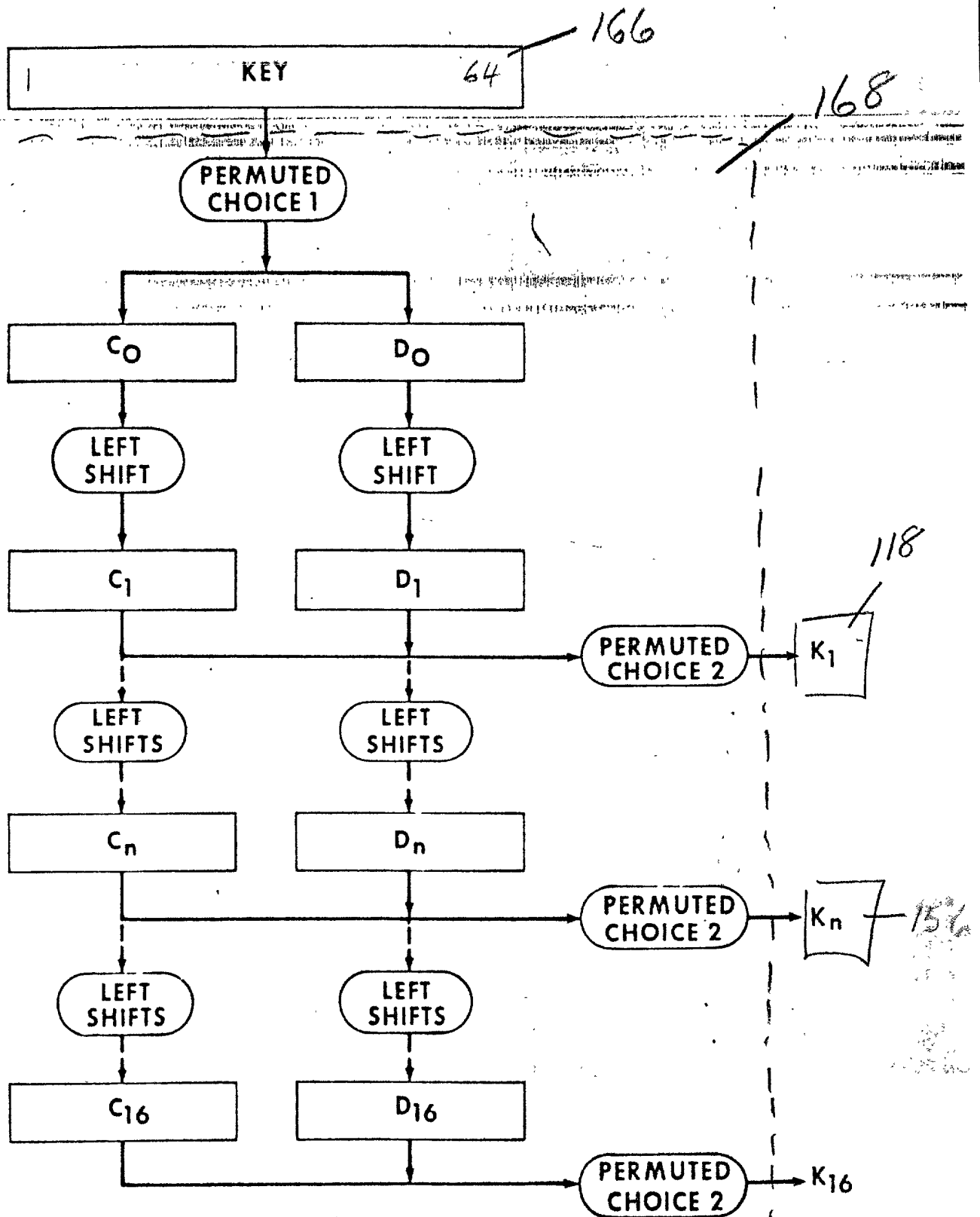


Fig. 7



Key schedule calculation.

Fig 8